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# FINITELY GENERATED SEMIGROUPS PRESENTED BY FINITE CONGRUENCE CLASSES II

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In this paper, we give a necessary and sufficient condition for one relator semigroups to be presented with finite congruence classes in the case of one relator of a special form. under an assumption .

## 1 Finitely generated monoid and their presentations

**Definition** Let  $X$  be a finite set of alphabets and  $R$  a finite subset of  $X^* \times X^*$ . Then  $R$  is *string-rewriting system*. Define the reduction relation  $\Rightarrow_R$  on  $X^*$  by  $\Rightarrow_R = \{((uw_1v, uw_2v) | u, v \in X^*, (w_1, w_2) \in R)\}$ . For  $u, v \in X^*$ ,  $(w_1, w_2) \in R$ , use the denotation :  $uw_1v \Rightarrow_R uw_2v$ . The congruence  $\mu_R$  on  $X^*$  (or  $X^+$ ) generated by  $\Rightarrow_R$  is called the *Thue congruence* defined by  $R$ . A monoid  $S$  has a *finite presentation* if there exists a finite set of  $X$ , there exists a surjective homomorphism  $\phi$  of  $X^*$  to  $S$  and there exists a string-rewriting system  $R$  consisting of pairs of words over  $X$  such that the Thue congruence  $\mu_R$  is the congruence  $\{(w_1, w_2) \in X^* \times X^* | \phi(w_1) = \phi(w_2)\}$ . Further, if for each  $w \in X^*$ , the congruence classes  $\mu_R(w) = \{w' \in X^* | (w, w') \in \mu_R\}$  is finite, then the monoid  $S = X^*/\mu_R$  is called to be *presented by finite congruence classes*. (Refer to [2],[3] and see [1] for examples)

If  $R = \{(u, v)\}$  then we say that  $R$  is an *one relator* and  $S$  is an *one relator monoid*.

## 2 The main theorems

First we have

**Theorem 1.** *Let  $u, v$  be word over a finite alphabet  $X$  and  $R = \{(u, w)\}$  a one-relator rewriting system. Assume that  $u$  is an unbordered and the length of  $u$  is shorter than one of  $v$  . Further,*

assume that  $u$  is not a subword of  $v$  and  $v$  contain at least one letter which  $u$  does not contain. Then the relator  $R = \{(u, v)\}$  does not generate the congruence such that all of the congruence classes are finite if and only if there exist non-empty words  $l_{i,j}, r_{i,j}$  over  $X$  such that  $u = l_{s,t}r_{s,t} (1 \leq s \leq 2k, 1 \leq t \leq i_s), u = l_0r_0,$

$$v \in X^+l_{1,i_1} \cdots l_{1,1}l_0, v \in r_{1,i_1-1}X^+ \cap \cdots \cap r_{1,1}X^+,$$

$$v \in r_{1,i_1}r_{2,1} \cdots r_{2,i_2}X^+, \quad v \in X^+l_{2,1} \cap \cdots \cap X^+l_{2,i_2-1}$$

$$v \in X^+l_{2k+1,i_{2k+1}} \cdots l_{2k+1,1}l_{2k,i_{2k}}, v \in r_{2k+1,i_{2k+1}-1}X^+ \cap \cdots \cap r_{2k+1,1}X^+,$$

$$\text{and } l_{2k+1,i_{2k+1}} = l_{1,i_1}.$$

Then Theorem 2 follows from Theorem 1.

**Theorem 2.** *Under the same assumption, the problem of whether one relator monoid  $S = X^* / \langle (u, v) \rangle$  are presented by finite congruence classes or not is decidable.*

## References

- [1] P.M. Higgins, *Techniques of semigroup theory*, Oxford Univ. press, 1992.
- [2] K. Shoji, *Finitely generated semigroups which have such a presentation that all the congruence classes are regular language*, Math. Japonica, **69**(2008), 73-78.
- [3] K. Shoji, *Finitely generated semigroups presented by finite congruence classes*, Surikaiseikikennkyuujo kokyuroku **1809**(2012), 160-170.